

Quantum Computing: Recap

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Contents

- Quantum Mechanics Refresher
 - states, operators, evolution, measurements
 - composite quantum systems, direct product structure, entanglement
- Quantum Algorithms for perfect Quantum Computers
 - Deutsch's problem and Simon's problem
 - quantum circuits (teleportation)
 - quantum computers can do any classical computation
 - simulating quantum computers is exceedingly hard
 - Quantum Fourier transform
 - phase estimation

Contents

- Quantum Algorithms for perfect Quantum Computers
 - Quantum Fourier transform, phase estimation
 - Shor's algorithm
 - HHL algorithm (solving systems of linear equations)
- Quantum Error Correction
 - repetition codes stabilizer formalism, surface codes
- NISQ Quantum Computing
 - variational algorithms
 - QAOA
 - Cirq

Composite Quantum Systems

superposition of states is a possible state
for qubits:

2 $|\psi\rangle = c_{00} |00\rangle + c_{01} |01\rangle + c_{10} |10\rangle + c_{11} |11\rangle$

3 $|\psi\rangle = c_{000} |000\rangle + c_{001} |001\rangle + c_{010} |010\rangle + c_{011} |011\rangle$
 $= c_{100} |100\rangle + c_{101} |101\rangle + c_{110} |110\rangle + c_{111} |111\rangle$

4 $|\psi\rangle = c_{0000} |0000\rangle + c_{0001} |0001\rangle + c_{0010} |0010\rangle + c_{0011} |0011\rangle$
 $= c_{0100} |0100\rangle + c_{0101} |0101\rangle + c_{0110} |0110\rangle + c_{0111} |0111\rangle$
 $= c_{1000} |1000\rangle + c_{1001} |1001\rangle + c_{1010} |1010\rangle + c_{1011} |1011\rangle$
 $= c_{1100} |1100\rangle + c_{1101} |1101\rangle + c_{1110} |1110\rangle + c_{1111} |1111\rangle$

Composite Quantum Systems

entanglement is important

$$\begin{aligned} |\psi\rangle &= (a_0 |0\rangle + a_1 |1\rangle) \otimes (b_0 |0\rangle + b_1 |1\rangle) \\ &= a_0 b_0 |00\rangle + a_0 b_1 |01\rangle + a_1 b_0 |10\rangle + a_1 b_1 |11\rangle \end{aligned}$$

$$|\psi\rangle = c_{00} |00\rangle + c_{01} |01\rangle + c_{10} |10\rangle + c_{11} |11\rangle$$

could express exponential amount of number with much less numbers

reduced density matrix of (maximally) entangled state

$$\langle 0| |\phi\rangle \langle \phi| |0\rangle + \langle 1| |\phi\rangle \langle \phi| |1\rangle = \frac{1}{2} (|0\rangle \langle 0| + |1\rangle \langle 1|)$$

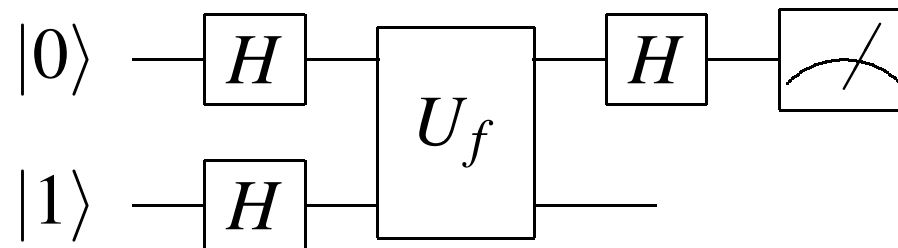
no local information

Quantum Algorithms

in this part assumed:

- gates work perfectly
- qubits stay in the same state if not acted upon
→ infinite coherence times

example:



often initially
Hadamards on all qubits

$$H^n |x\rangle = \frac{1}{2^{n/2}} \sum_{y \in \{0,1\}^n} (-1)^{x \cdot y} |y\rangle$$

$$x = x_0, x_1, \dots, x_n$$

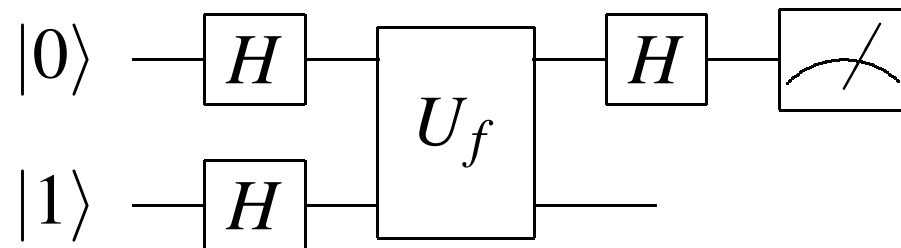
superposition of all
basis states

Quantum Algorithms

in this part assumed:

- gates work perfectly
- qubits stay in the same state if not acted upon
→ infinite coherence times

example:



make sure no correlations
between qubits before readout

for maximally entangled state:

$$\langle 0| |\phi\rangle \langle \phi| |0\rangle + \langle 1| |\phi\rangle \langle \phi| |1\rangle = \frac{1}{2} (|0\rangle \langle 0| + |1\rangle \langle 1|)$$

measurement (readout) produces random result

Quantum Algorithms

quantum evolution unitary \rightarrow algorithms reversible

$$|\psi_f\rangle = U |\Psi_i\rangle$$

\rightarrow decompose U into single- and two-qubit gates

\rightarrow is the only stuff that can be implemented

\rightarrow entangling two-qubit gates are important, e.g. CNOT

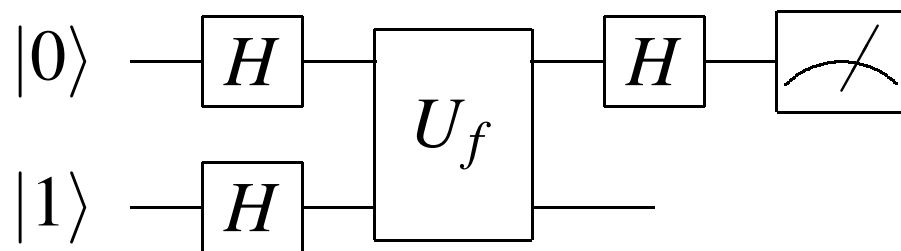
$$(\alpha |0\rangle + \beta |1\rangle) \otimes |0\rangle \rightarrow \alpha |00\rangle + \beta |11\rangle$$

arbitrary single-qubit gates plus CNOT is universal

\rightarrow can do any computation

Deutsch's and Simon's Problems

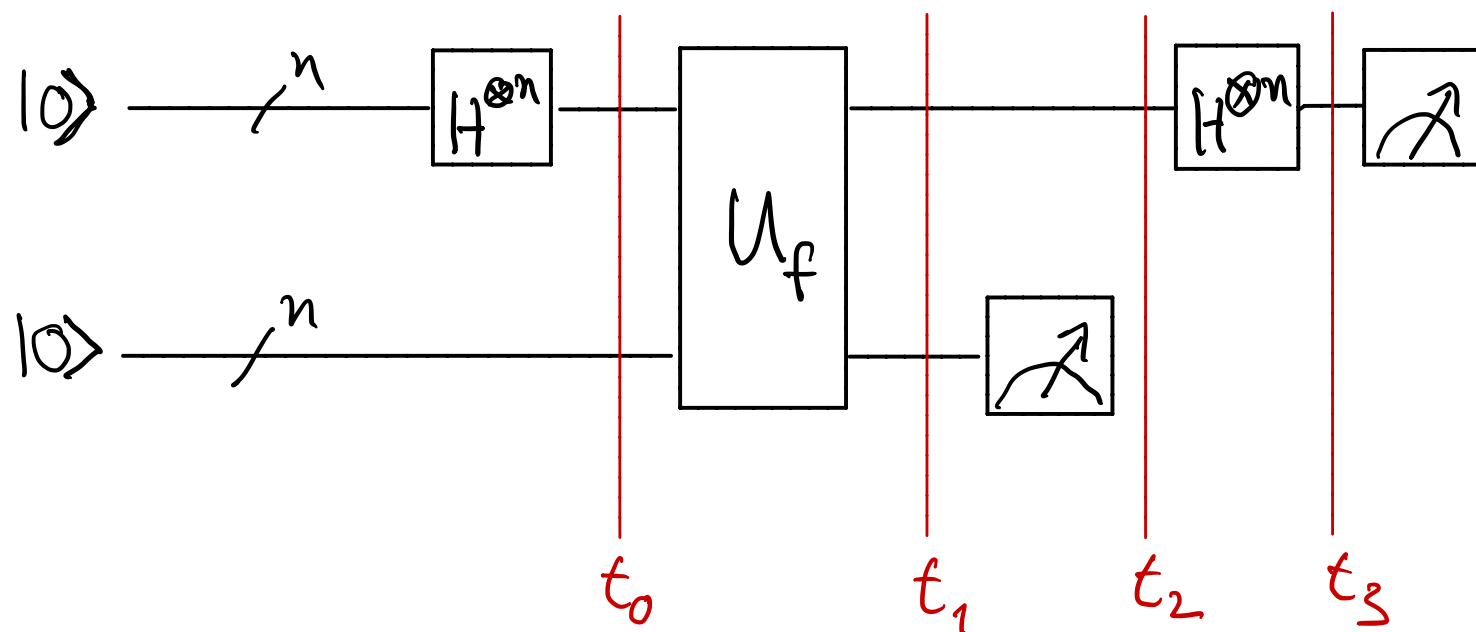
Deutsch: $f(0) = f(1)$?



$$U_f |j\rangle |k\rangle = |j\rangle |f(j) \oplus k\rangle$$

class.: 2 queries
quantum: 1 query

Simon: find period s of function ($f(x) = f(y)$ for $y = x \oplus s$)

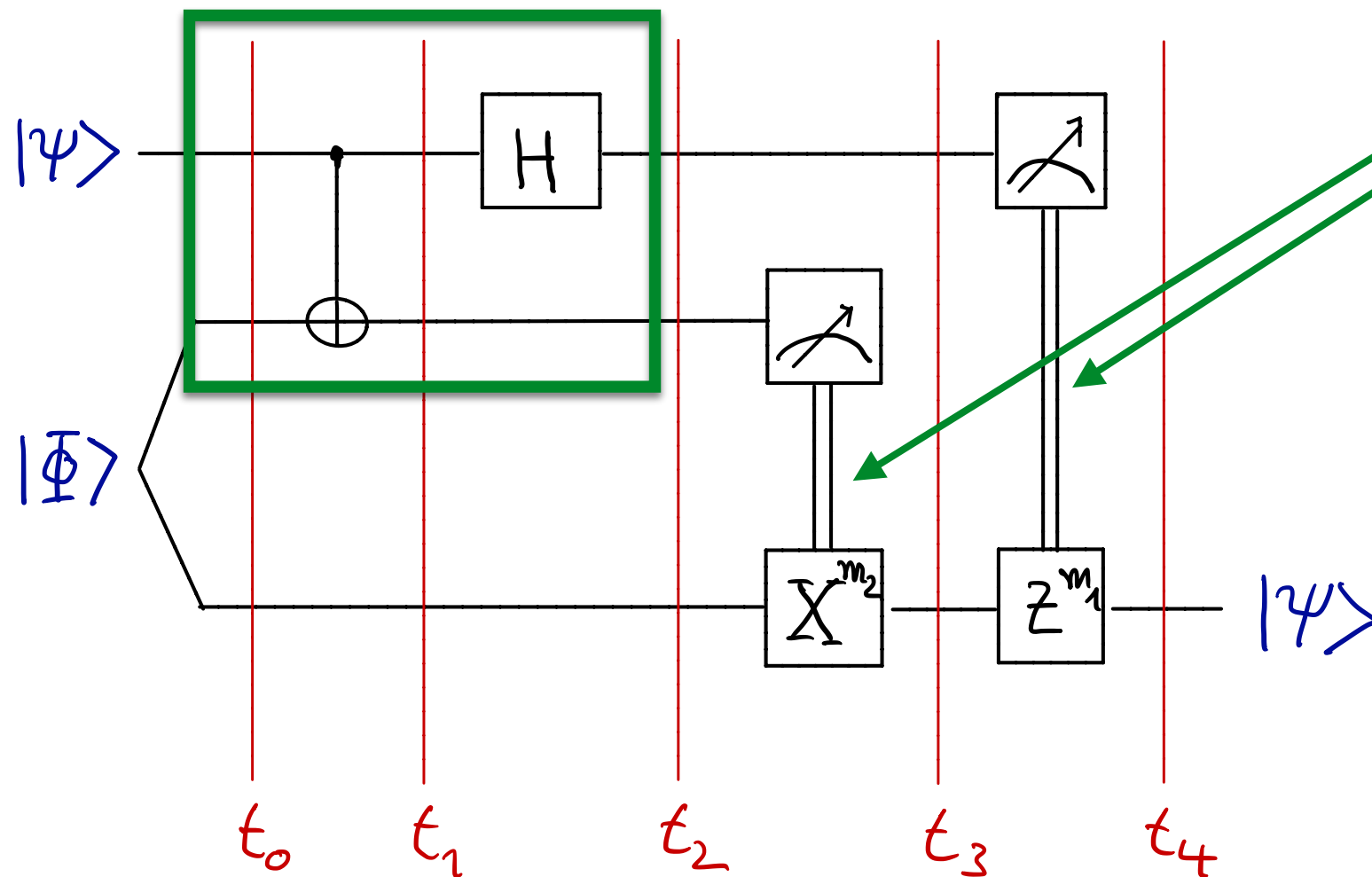


$$U_f |j\rangle |k\rangle = |j\rangle |f(j) \oplus k\rangle$$

class.: $2^{n/2}$ queries
quantum: $\Omega(n)$ queries

Quantum Teleportation

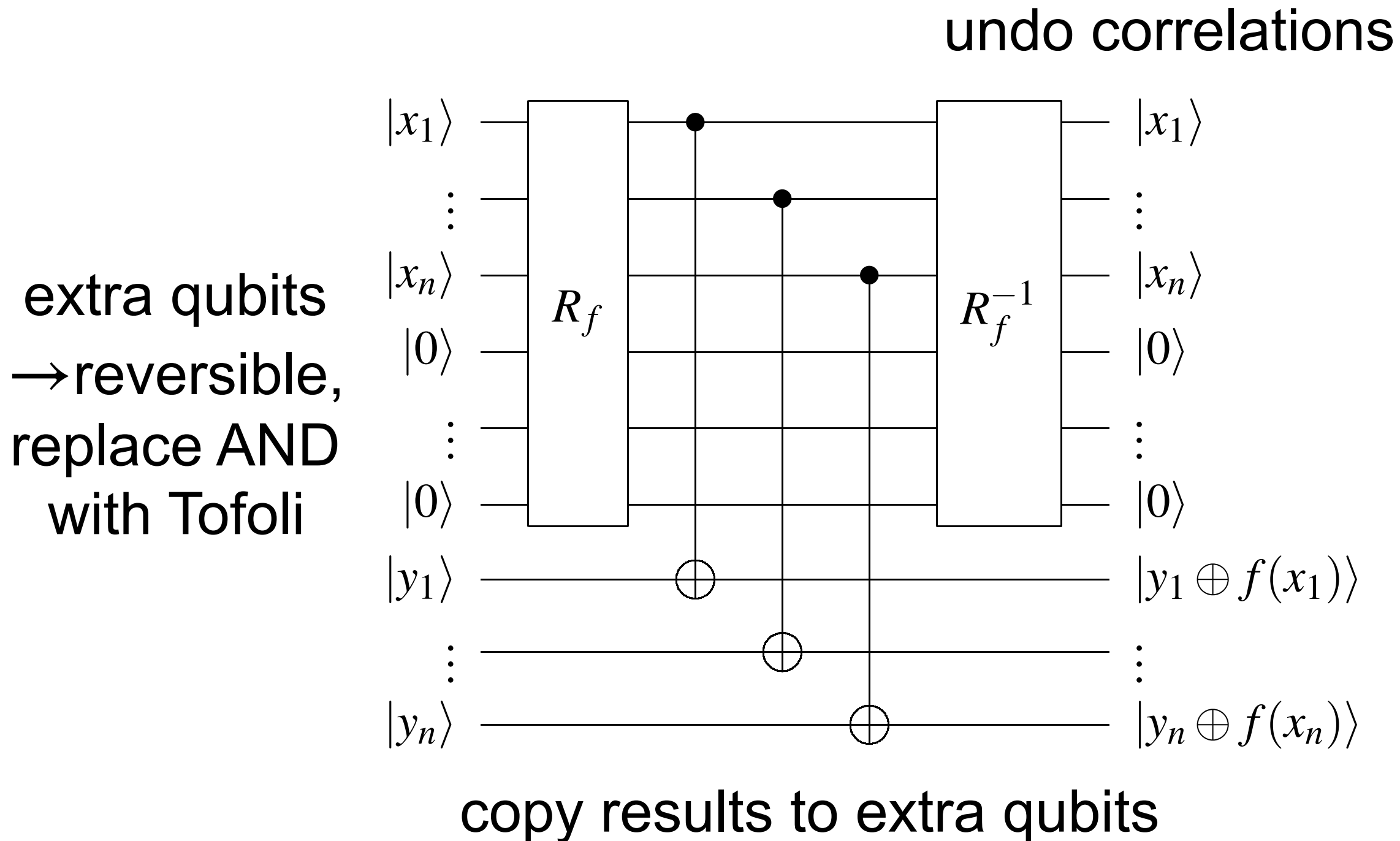
state transfer circuit



transfer does not
happen before
classical
information is
transmitted!

Classical vs Quantum

any classical algorithm can be implemented
on a quantum computer



Classical vs Quantum

simulating a quantum computer requires
exponential resources

Schrödinger

$$|\psi\rangle = \sum_{j_1, j_2, \dots, j_n=0}^1 \alpha_{j_1, j_2, \dots, j_n} |j_1, j_2, \dots, j_n\rangle = \sum_{\vec{j}} \alpha_{\vec{j}} |\vec{j}\rangle$$

$$\frac{d}{dt} \alpha_{\vec{j}} = \sum_{\vec{l}} H_{\vec{j}; \vec{l}} \alpha_{\vec{l}}$$

exponential
amount of RAM

Feynman

$$\begin{aligned} \langle \vec{j}_f | U | \vec{j}_i \rangle &= \langle \vec{j}_f | U_L U_{L-1} \dots U_2 U_1 | \vec{j}_i \rangle \\ &= \sum_{\vec{j}_{L-1}} \sum_{\vec{j}_{L-2}} \dots \sum_{\vec{j}_1} \langle \vec{j}_f | U_L | \vec{j}_{L-1} \rangle \langle \vec{j}_{L-1} | U_{L-1} | \vec{j}_{L-2} \rangle \dots \langle \vec{j}_1 | U_1 | \vec{j}_i \rangle \end{aligned}$$

exponential amount of time

Quantum Fourier Transform

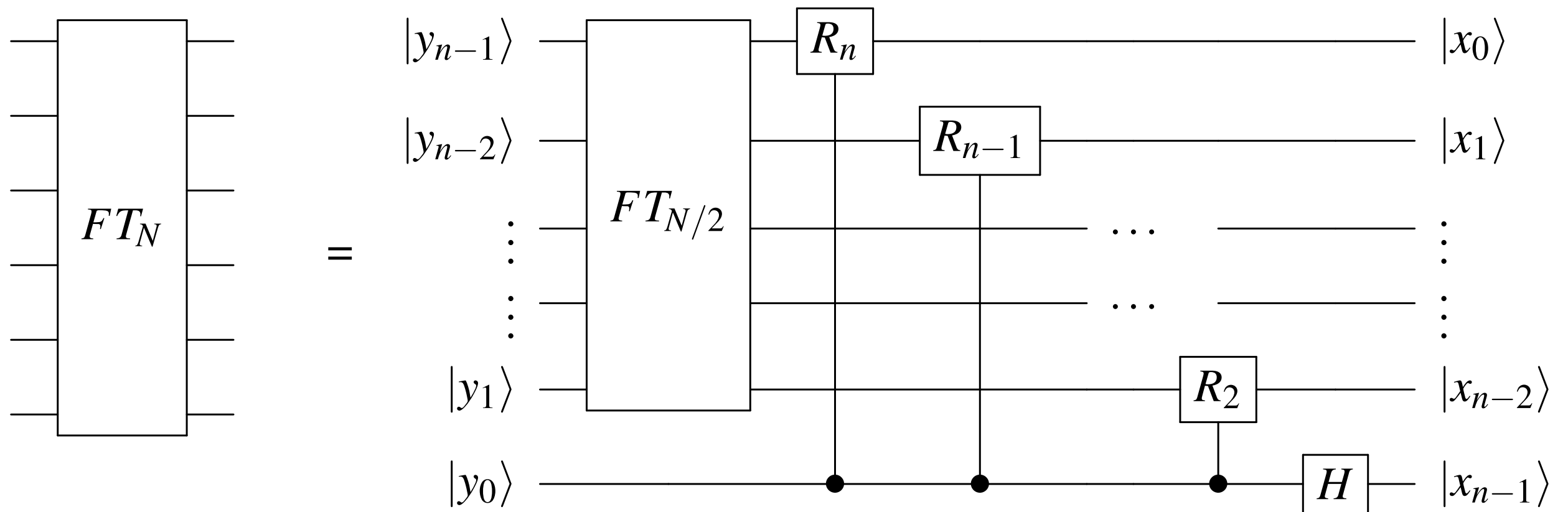
$$|\tilde{y}\rangle = \frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} \omega^{xy} |x\rangle$$

$$\omega = e^{2\pi i/N}$$

$$\omega^N = 1$$

$$\sum_{j=0}^{N-1} \omega^j = 0$$

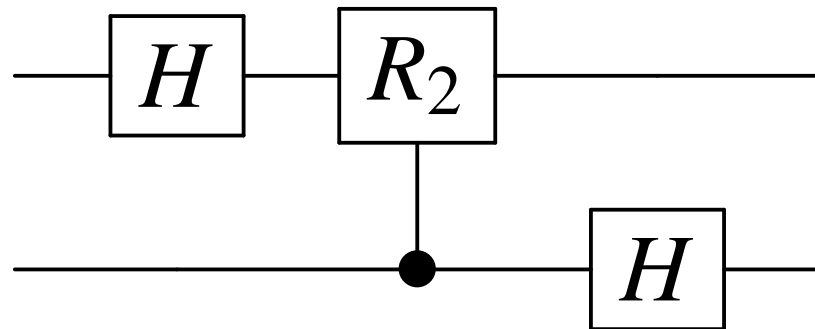
$$|x\rangle = |x_{n-1}, x_{n-2}, \dots, x_0\rangle \quad x = x_{n-1}2^{n-1} + x_{n-2}2^{n-2} + \dots + x_12 + x_0$$



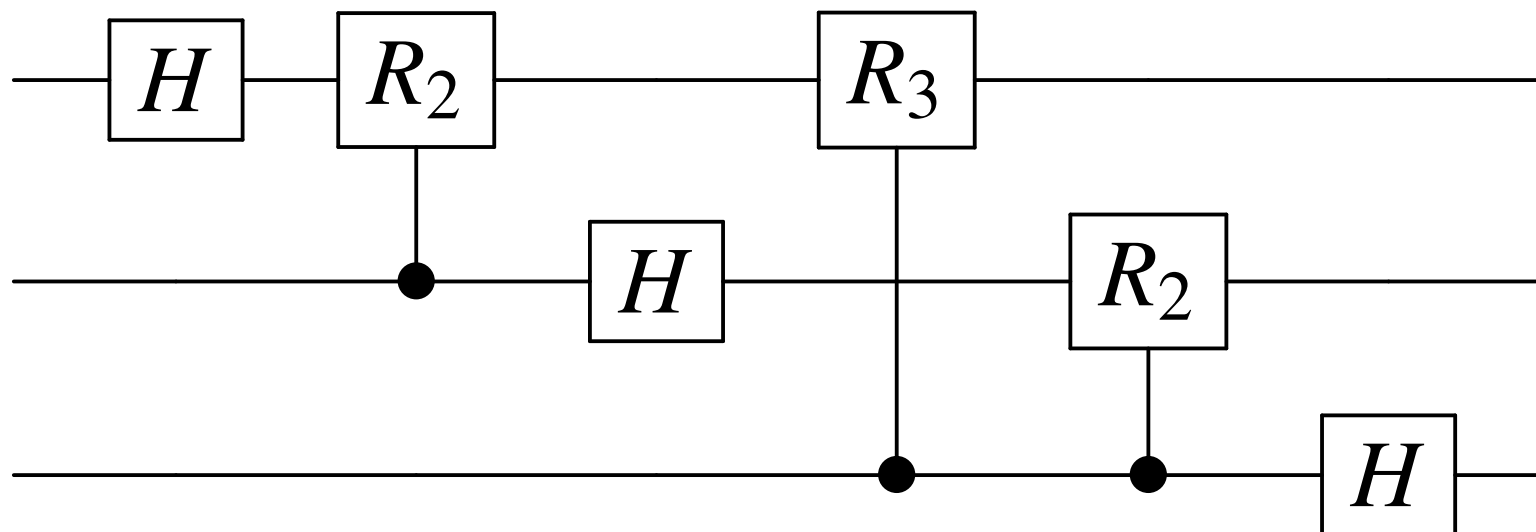
$$R_j = |0_j\rangle\langle 0_j| + e^{2\pi i 2^{-j}} |1_j\rangle\langle 1_j|$$

Quantum Fourier Transform

$n=2$:



$n=3$:



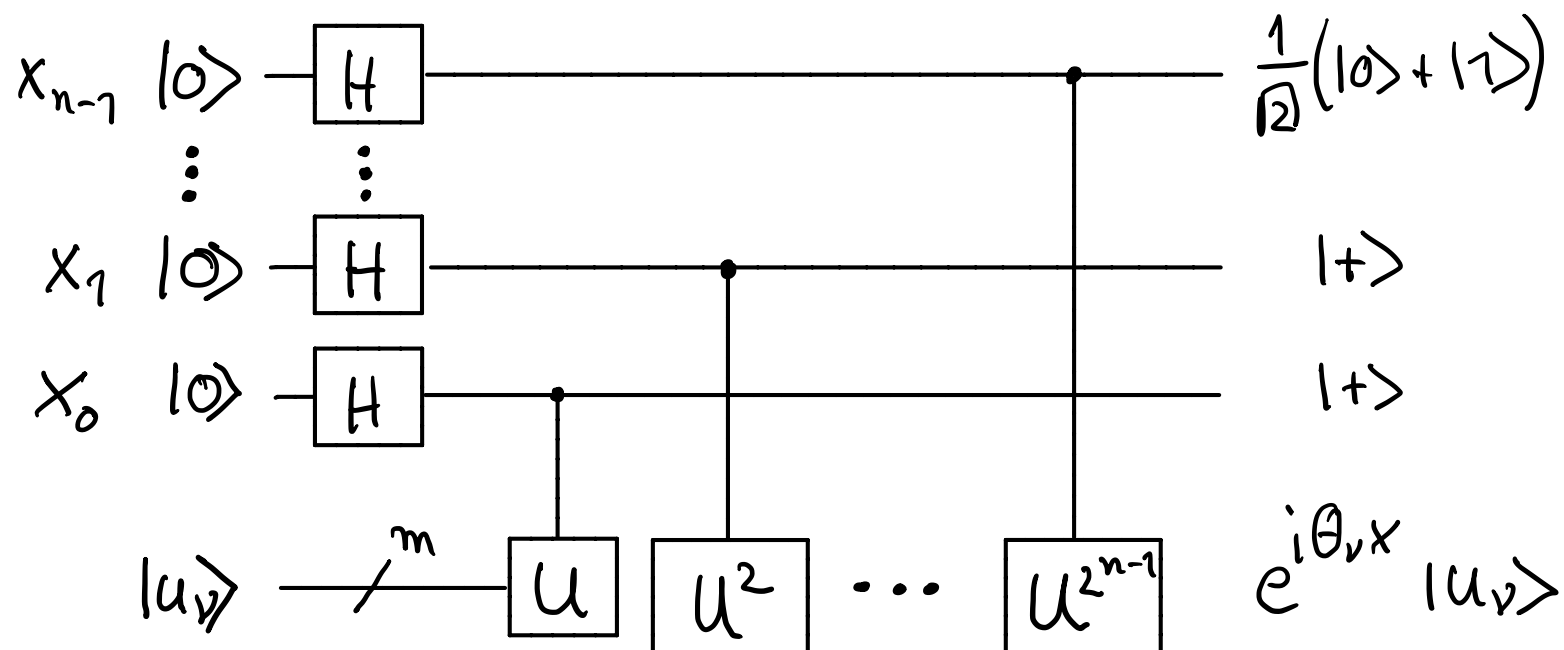
Quantum Phase Estimation

$$U |u_v\rangle = e^{i\theta_v} |u_v\rangle \quad \text{find } \theta_v$$

$$|x\rangle = |x_{n-1}, x_{n-2}, \dots, x_0\rangle \quad x = x_{n-1}2^{n-1} + x_{n-2}2^{n-2} + \dots + x_12 + x_0$$

$$FT_N^\dagger \frac{1}{2^{n/2}} \sum_{x=0}^{N-1} e^{i\theta_v x} |x\rangle = \sum_y \left(\frac{1}{N} \sum_x \exp \left[\frac{2\pi i}{N} \left(\frac{N\theta_v}{2\pi} - y \right) x \right] \right) |y\rangle$$

prepare via



strongly peaked at
 $y = N\theta_v/(2\pi)$
 \rightarrow will measure
 $y = N\theta_v/(2\pi)$

Shor's Algorithm

quantum step: period finding

find period of $f(x) = a^x \bmod N$ ($a > 1, x > 0$)

$U |x\rangle |b \bmod N\rangle = |x\rangle |b a \bmod N\rangle$ controlled multiplication

$U^x |x\rangle |b \bmod N\rangle = |x\rangle |b a^x \bmod N\rangle$ as in phase estimation

→ apply phase estimation to U $U |u_v\rangle = e^{i\theta_v} |u_v\rangle$

can prepare $|1\rangle = |a^0\rangle$: superposition of all eigenstates of U

multiple phase estimation measurements → period of $f(x)$

HHL Algorithm

solve: $A |x\rangle = |b\rangle \rightarrow |x\rangle \propto A^{-1} |b\rangle$

strategy: phase estimation for $U = \exp(2\pi i \kappa A / N)$

apply controlled rotation of form:

$$|y\rangle |0\rangle \rightarrow |y\rangle \left(\sqrt{1 - \frac{y^2}{2}} |0\rangle + \frac{1}{y} |1\rangle \right)$$

eigenstate of U

measure and keep outcome 1

\rightarrow applied $1/\lambda_a$, where $A |a\rangle = \lambda_a |a\rangle$

Error Correction

gates will never be perfect \rightarrow actively correct for errors

many physical qubits form a logical qubit

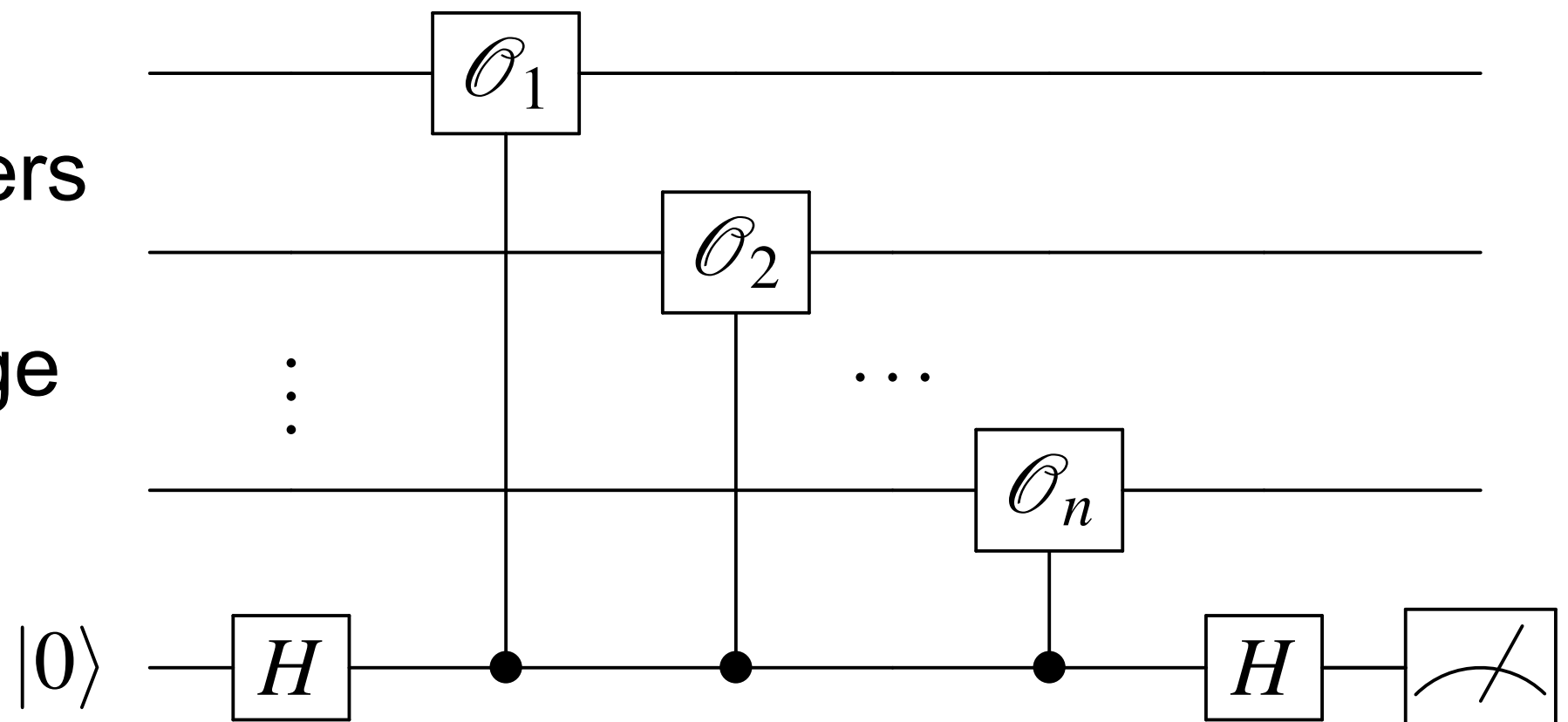
logical states \rightarrow entangled states of many physical qubits

\rightarrow stabilizers

$$P_k |\psi_j\rangle_L = |\psi_j\rangle_L$$

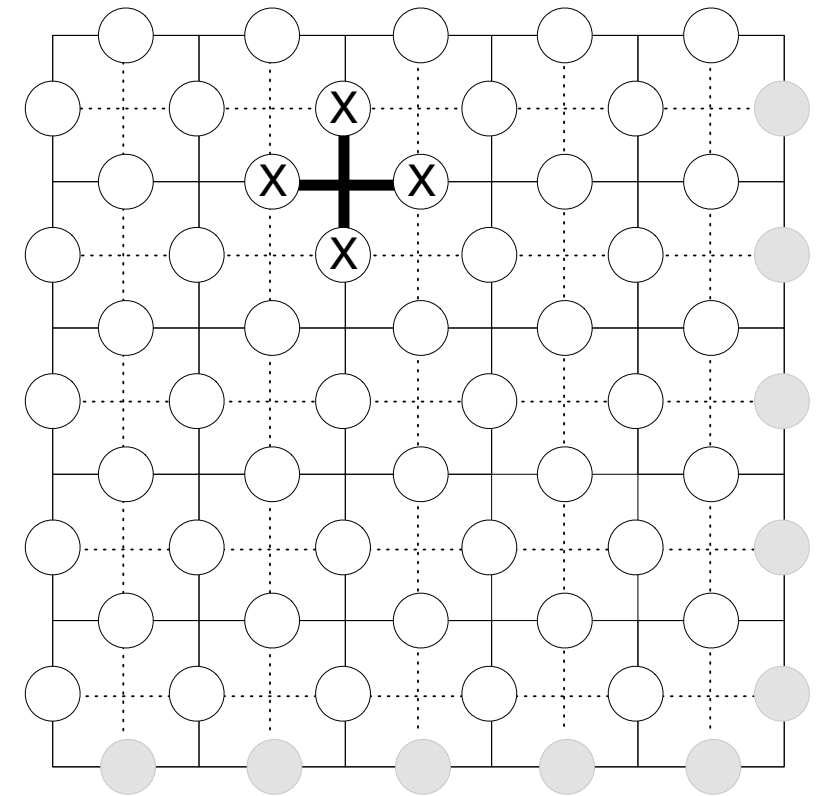
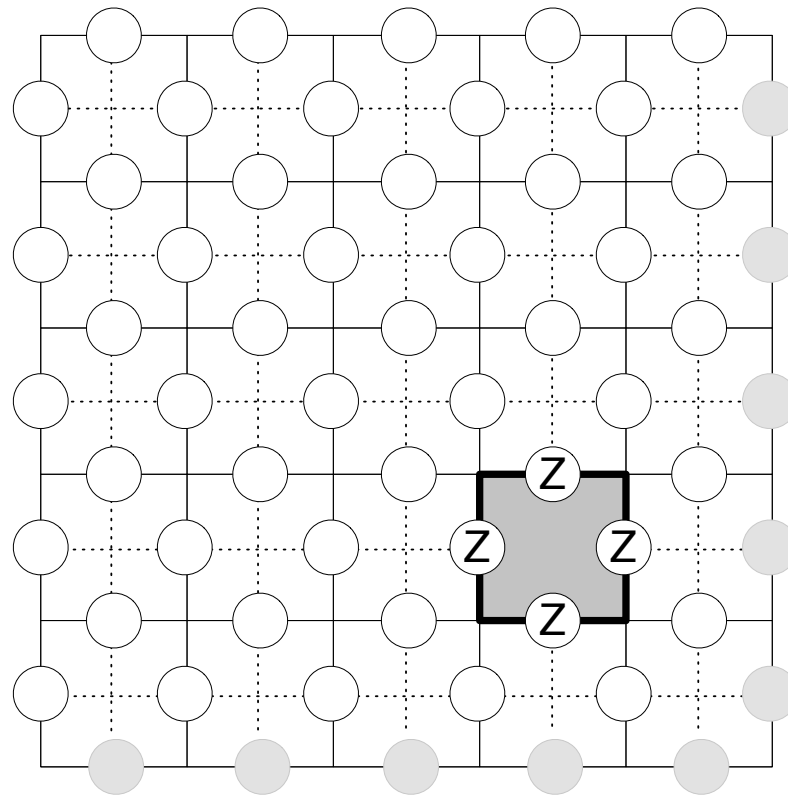
$$P_k = X \otimes X \otimes Y \otimes Z \otimes \mathbb{1} \otimes \dots$$

measuring stabilizers
detects errors
but does not change
logical states



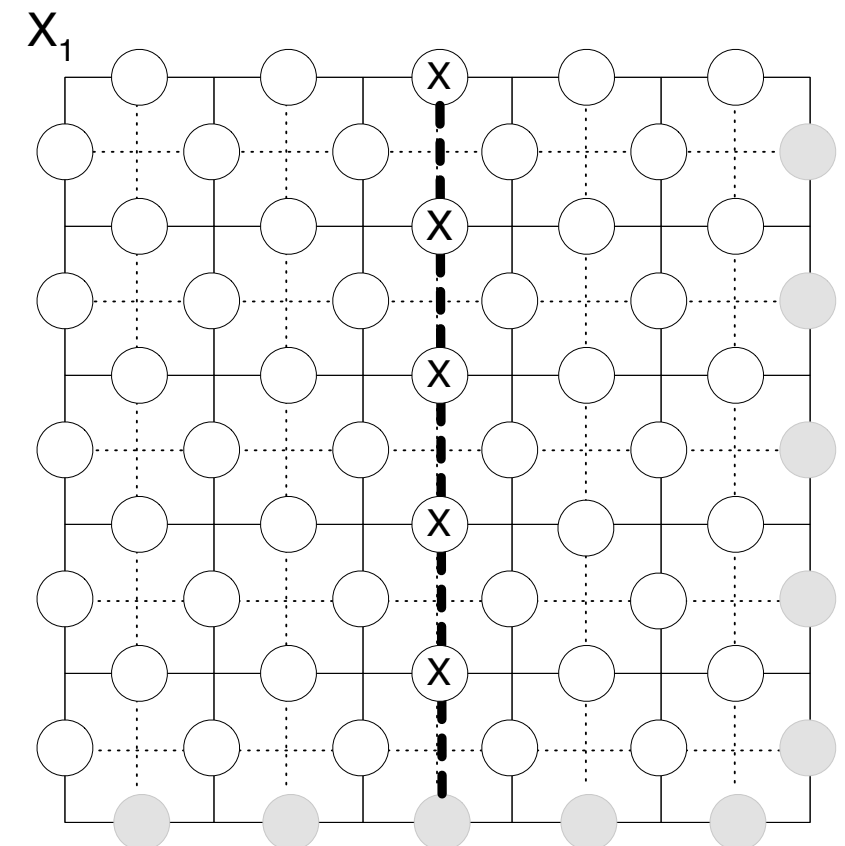
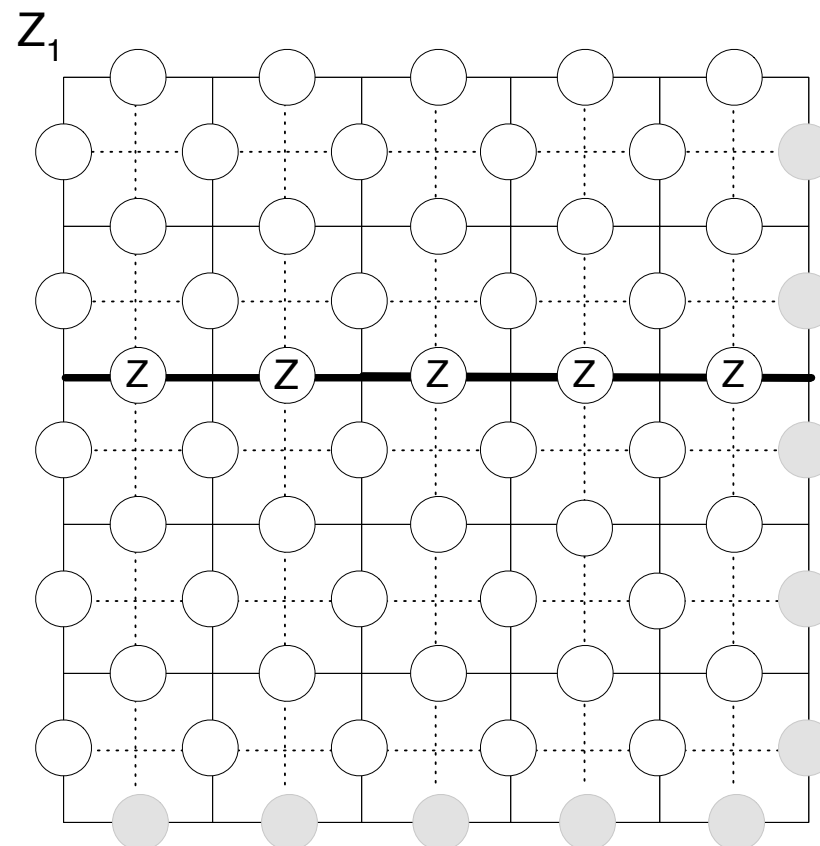
Error Correction: Toric Code

stabilizer
generators



4 logical states
2 logical qubits

logical
operators



Error Correction: Surface Code

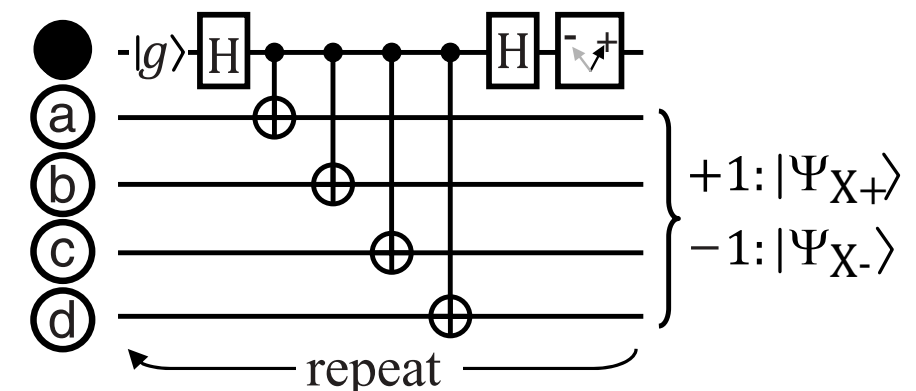
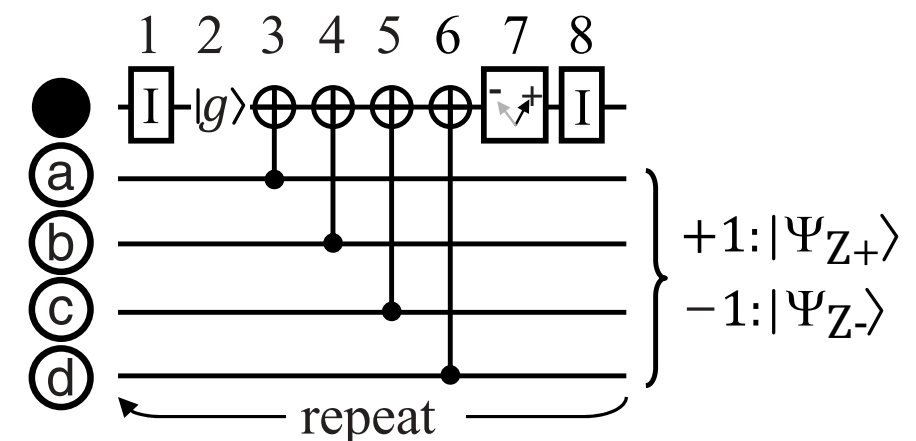
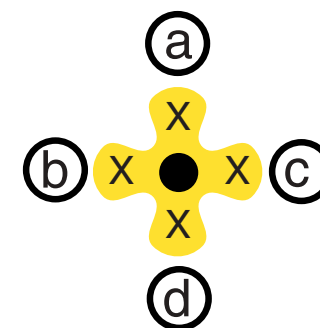
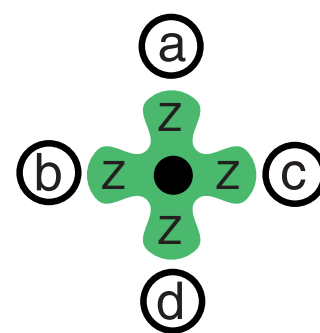
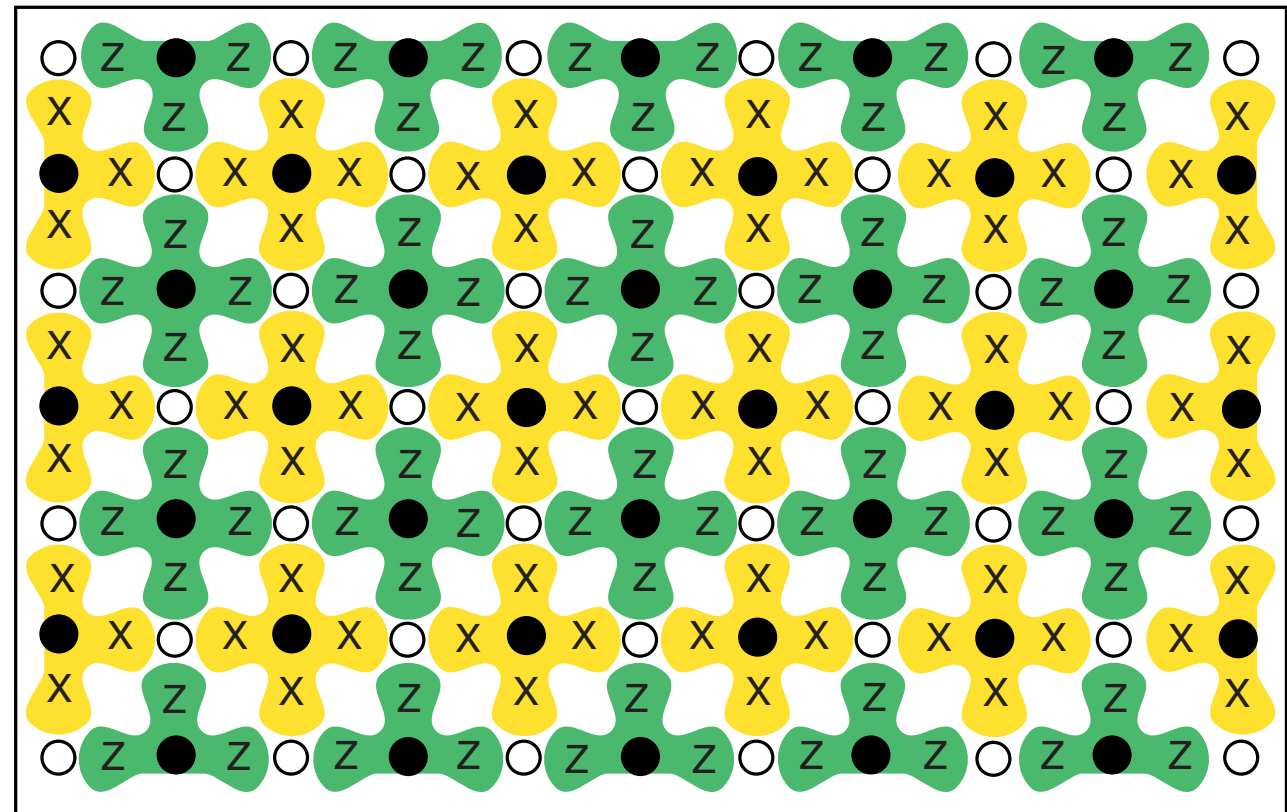
real hardware

→ no periodic
boundary
condition

errors can also occur
during error detection
and correction

→ need $p_{error} < 1\%$

→ need 1000 physical
qubits per logical qubit



Variational Algorithms

will not have error correction soon → what can we do now?

find ground state of a Hamiltonian H

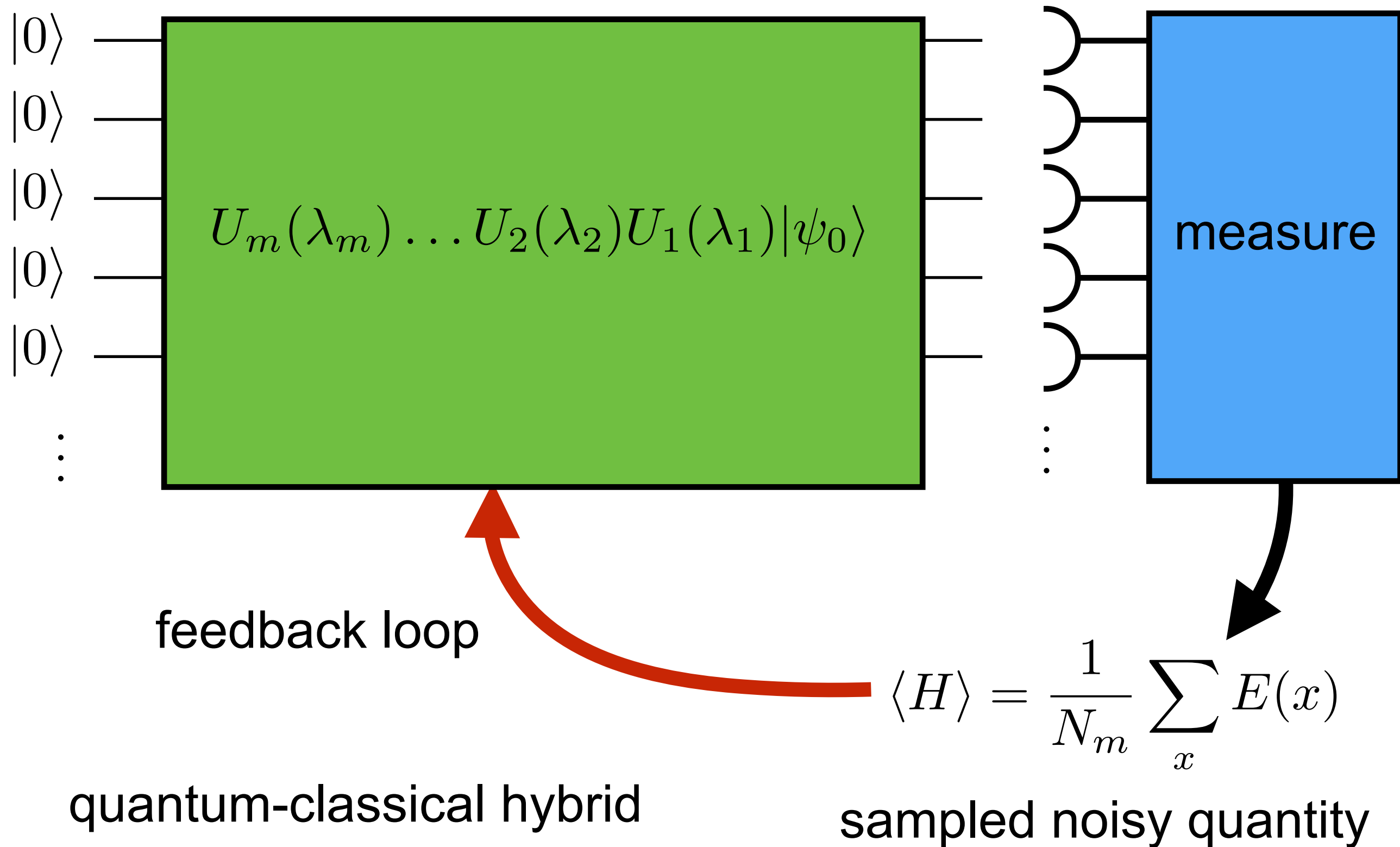
- condensed matter physics
- ground or low energy state
 - stable configuration
 - quantum chemistry
- can be applied to combinatorial optimization
 - QAOA

$$|\psi(\vec{\theta})\rangle = U_m(\theta_m) U_{m-1}(\theta_{m-1}) \dots U_1(\theta_1) |0, 0, \dots, 0\rangle$$

$\theta_1, \theta_2, \dots, \theta_m$ are variational parameters

$$\rightarrow \text{minimize } E(\vec{\theta}) = \langle \psi(\vec{\theta}) | H | \psi(\vec{\theta}) \rangle$$

Variational Algorithms



Quantum Approximate Optimization

combinatorial optimization problem:

find bit-string $z = z_1 z_2 \dots z_n$ such that $C(z)$ is minimized

sum of clauses: $C(z) = -C_1(z) - C_2(z) - C_3(z) - \dots$

all clauses true, $C_j(z) = 1$, \rightarrow minimum of $C(z)$

example

$$C(z) = 3z_1z_2 - z_2z_3 + z_4$$



make
quantum

$$C(Z) = 3Z_1Z_2 - Z_2Z_3 + Z_4$$

\rightarrow Hamiltonian \rightarrow variational algorithm

Quantum Approximate Optimization

search $|\vec{z}\rangle = |z_1, z_2, \dots, z_n\rangle$ such that $C(Z)|\vec{z}\rangle = E_0|\vec{z}\rangle$

ansatz:

$$|\psi(\vec{\gamma}, \vec{\beta})\rangle = U_B(\beta_p)U_C(\gamma_p)U_B(\beta_{p-1})U_C(\gamma_{p-1}) \dots U_B(\beta_1)U_C(\gamma_1)|s\rangle$$

$$U_C(\gamma) = e^{-i\gamma C(Z)}$$

$$U_B(\beta) = e^{-i\beta \sum_j X_j}$$

$$|s\rangle = H^{\otimes n} |0, 0, \dots, 0\rangle = \frac{1}{\sqrt{2^n}} \sum_z |z\rangle$$

variational algorithm

$$\rightarrow \gamma_1, \gamma_2, \dots, \gamma_p$$

$$\beta_1, \beta_2, \dots, \beta_p$$

Quantum Approximate Optimization

search $|\vec{z}\rangle = |z_1, z_2, \dots, z_n\rangle$ such that $C(Z) |\vec{z}\rangle = E_0 |\vec{z}\rangle$

we know that solution is a product state

variational algorithm \rightarrow

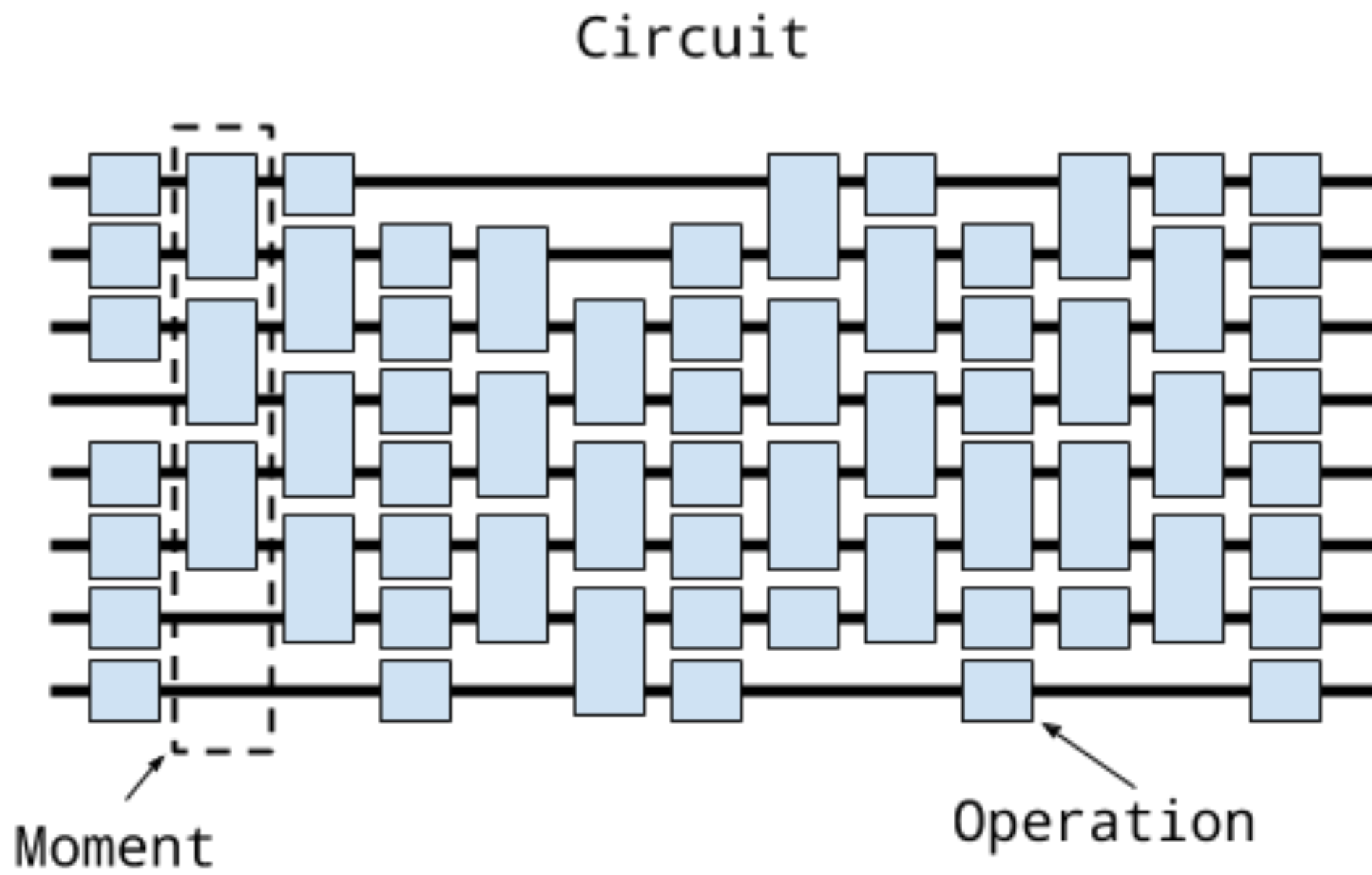
$$|\psi(\vec{\gamma}, \vec{\beta})\rangle = \sum_{z_1, z_2, \dots, z_n=0}^1 \alpha_{z_1, z_2, \dots, z_n} |z_1, z_2, \dots, z_n\rangle$$

will measure z_1, z_2, \dots, z_n with probability $|\alpha_{z_1, z_2, \dots, z_n}|^2$

for each measured z_1, z_2, \dots, z_n can compute $C(z)$

can find solution as long as it contributes to $|\psi(\vec{\gamma}, \vec{\beta})\rangle$

Cirq



Thanks!

And all the best for
the exam!